

# Predicate Logic

## Translation

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PHIL 500

$$\forall x((Px \vee Qx) \rightarrow Rx)$$

# Outline

Four Important Statement Forms

2-Place Predicates

Syntax for PL

Vocabulary

Grammar

Free and Bound Variables

Important Syntactic Features in PL

# Four Important Statement Forms

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(A) All  $\mathcal{F}$ s (in the domain) are  $\mathcal{G}$ s

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## Variants of 'All $\mathcal{F}$ s are $\mathcal{G}$ s'

- ▷ All  $\mathcal{F}$ s are  $\mathcal{G}$ s
- ▷ Any  $\mathcal{F}$  is a  $\mathcal{G}$
- ▷ Every  $\mathcal{F}$  is  $\mathcal{G}$

## Variants of 'No $\mathcal{F}$ s are $\mathcal{G}$ s'

- ▶ No  $\mathcal{F}$ s are  $\mathcal{G}$ s
- ▶ No  $\mathcal{F}$  is  $\mathcal{G}$
- ▶ No  $\mathcal{F}$  is a  $\mathcal{G}$
- ▶ There are no  $\mathcal{G}$   $\mathcal{F}$ s

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- ▶ Some  $\mathcal{F}$ s are not  $\mathcal{G}$ s
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- ▶ Some  $\mathcal{F}$  is not  $\mathcal{G}$
- ▶ Some  $\mathcal{F}$  is a non- $\mathcal{G}$
- ▶ There are non- $\mathcal{G}$   $\mathcal{F}$ s

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- Some of you may find one or more of these alternative translations more natural—if so, you should feel free to use them instead.

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- ▶ If you don't find any of these more natural, don't worry about it—just ignore this slide.

# Examples

domain : all people

*F*\_\_ : \_\_ is funny

*S*\_\_ : \_\_ is shy

*T*\_\_ : \_\_ is tall

*Q*\_\_ : \_\_ is quirky

Everyone is funny :

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Everyone is funny :

$\forall x Fx$

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Someone is quirky :

$\exists y Qy$

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Everyone tall is shy :

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Everyone tall is shy :

$$\forall z (Tz \rightarrow Sz)$$

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$$\forall x [Sx \rightarrow Qx] \rightarrow \forall x [Sx \rightarrow Fx]$$

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Some tall people are shy :

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Some tall people are shy :

$$\exists w (T w \wedge S w)$$

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$$\exists w (Tw \wedge Sw)$$

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No tall people are either funny or quirky :

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$$\forall x [Tx \rightarrow \neg(Fx \vee Qx)]$$

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Some tall people are neither funny nor shy :

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*V* \_\_\_ : \_\_\_ is vegetarian

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*a* : Albert

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Some people are vegetarian :

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Some people are vegetarian :

$$\exists x(\mathcal{F}x \wedge \mathcal{G}x)$$

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$\exists x Vx$

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Some ferocious animals are not carnivorous :

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Some people are vegetarians and some are not :

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$$\exists x(Px \wedge Vx) \wedge \exists y(Py \wedge \neg Vy)$$

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If Albert is ferocious, then all people are ferocious :

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$$Fa \rightarrow \forall x (Fx \rightarrow Gx)$$

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Albert is ferocious if anyone is :

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Albert is ferocious if anyone is :

$$\exists x(\mathcal{F}x \wedge \mathcal{G}x) \rightarrow Fa$$

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$V$  \_\_\_ : \_\_\_ is vegetarian

$C$  \_\_\_ : \_\_\_ is carnivorous

$D$  \_\_\_ : \_\_\_ is ferocious

$a$  : Albert

If everyone is vegetarian, then no one is carnivorous :

# Examples

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If everyone is vegetarian, then no one is carnivorous :

$$\forall x(\mathcal{F}x \rightarrow \mathcal{G}x) \rightarrow \forall y(\mathcal{F}y \rightarrow \neg \mathcal{G}y)$$



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There are non-vegetarian people if and only if someone is ferocious :

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domain : all foods

*J*\_\_\_ : \_\_\_ is a jellybean

*B*\_\_\_ : \_\_\_ is black

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All black jellybeans are delicious, but no red jellybean is :

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$$\forall x [Jx \rightarrow Bx] \wedge \forall y [Jy \rightarrow \neg Ry]$$

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$$\forall x[(Bx \wedge Jx) \rightarrow Dx] \wedge \forall y[Fy \rightarrow \neg Dy]$$



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If some red jellybeans are delicious, then all black jellybeans are delicious :

$$\exists x [R x \wedge D x] \rightarrow \forall y [B y \rightarrow D y]$$

# Examples

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If some red jellybeans are delicious, then all black jellybeans are delicious :

$$\exists x[(Rx \wedge Jx) \wedge \mathcal{G}x] \rightarrow \forall y[\mathcal{F}y \rightarrow \mathcal{G}y]$$

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## 2-Place Predicates

---

# Predicates

- a *predicate* is a *gappy statement*—it's a statement with a name (or names) missing.

Tammy loves Sammy.

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- ▶ If a predicate has a single gap, then we'll call it a *1-place* predicate

# Predicates

- a *predicate* is a *gappy statement*—it's a statement with a name (or names) missing.

\_\_\_\_\_ loves \_\_\_\_\_

- ▷ If a predicate has a single gap, then we'll call it a *1-place* predicate
- ▷ If a predicate has two gaps, then we'll call it a *2-place* predicate

# Predicates

- a *predicate* is a *gappy statement*—it's a statement with a name (or names) missing.

\_\_\_\_\_ loves \_\_\_\_\_

- ▶ If a predicate has a single gap, then we'll call it a *1-place* predicate
- ▶ If a predicate has two gaps, then we'll call it a *2-place* predicate
- ▶ If a predicate has  $N$  gaps, then we'll call it an  *$N$ -place* predicate

## 2-Place Predicates

domain : all people

$L$  \_\_\_ \_\_\_ : \_\_\_ loves \_\_\_

$a$  : Abelard

$h$  : Heloise

## 2-Place Predicates

domain : all people

$L$  \_\_\_ \_\_\_ : \_\_\_ loves \_\_\_

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise :



## 2-Place Predicates

domain : all people

$L$  \_\_\_ \_\_\_ : \_\_\_ loves \_\_\_

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise :

Heloise loves Abelard :

## 2-Place Predicates

domain : all people

$L$  \_\_\_ \_\_\_ : \_\_\_ loves \_\_\_

$a$  : Abelard

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Abelard loves Heloise :

Heloise loves Abelard :

- ▶ We need some way of saying *which* gap is which

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

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Abelard loves Heloise :

Heloise loves Abelard :

- ▶ We need some way of saying *which* gap is which

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

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Abelard loves Heloise :  $Lah$

Heloise loves Abelard :

- ▶ We need some way of saying *which* gap is which

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise :  $Lah$

Heloise loves Abelard :  $Lha$

- ▶ We need some way of saying *which* gap is which

## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

Everyone loves Abelard :

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

$h$  : Heloise

Everyone loves Abelard :

$\forall x Lxa$



## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

Someone loves Heloise :

## 2-Place Predicates

domain : all people

$Lxy$  :  $\_\_x$  loves  $\_\_y$

$a$  : Abelard

$h$  : Heloise

Someone loves Heloise :

$\exists z Lzh$

## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise if anyone does :

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise if anyone does :

$$\exists x Lxh \rightarrow Lah$$

## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

Everyone who loves Heloise loves Abelard, too. :

## 2-Place Predicates

domain : all people

$Lxy$  :  $\_\_x$  loves  $\_\_y$

$a$  : Abelard

$h$  : Heloise

Everyone who loves Heloise loves Abelard, too. :

$$\forall x (Fx \rightarrow Gx)$$

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

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## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

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Everyone who loves Heloise loves Abelard, too. :

$$\forall x (Lxh \rightarrow Lxa)$$



## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

$h$  : Heloise

Abelard loves himself :

## 2-Place Predicates

domain : all people

$Lxy$  : \_\_\_ $x$  loves \_\_\_ $y$

$a$  : Abelard

$h$  : Heloise

Abelard loves himself :

$Laa$

## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

Everyone loves themselves.

## 2-Place Predicates

domain : all people

$Lxy$  :  $\text{---}x$  loves  $\text{---}y$

$a$  : Abelard

$h$  : Heloise

Everyone loves themselves.

$\forall z Lzz$

## 2-Place Predicates

domain : all people

$Mxy$  : \_\_\_ $y$  loves \_\_\_ $x$

$a$  : Abelard

$h$  : Heloise

## 2-Place Predicates

domain : all people

$Mxy$  : \_\_\_ $y$  loves \_\_\_ $x$

$a$  : Abelard

$h$  : Heloise

Everyone loves Abelard :

## 2-Place Predicates

domain : all people

$Mxy$  : \_\_\_ $y$  loves \_\_\_ $x$

$a$  : Abelard

$h$  : Heloise

Everyone loves Abelard :

$\forall x M ax$

## 2-Place Predicates

domain : all people

$Mxy$  :  $\text{---}y$  loves  $\text{---}x$

$a$  : Abelard

$h$  : Heloise

Someone loves Heloise :



## 2-Place Predicates

domain : all people

$Mxy$  : \_\_\_ $y$  loves \_\_\_ $x$

$a$  : Abelard

$h$  : Heloise

Someone loves Heloise :

$\exists z Mhz$

## 2-Place Predicates

domain : all people

$Mxy$  :  $\text{---}y$  loves  $\text{---}x$

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise if anyone does :

## 2-Place Predicates

domain : all people

$Mxy$  : \_\_\_ $y$  loves \_\_\_ $x$

$a$  : Abelard

$h$  : Heloise

Abelard loves Heloise if anyone does :

$$\exists x Mhx \rightarrow Mha$$

## 2-Place Predicates

domain : everything in the office

$j$  : Jim

$m$  : Michael

$p$  : Pam

$s$  : Stanley

$Lxy$  : \_\_\_ $x$  likes \_\_\_ $y$

$Ex$  : \_\_\_ $x$  is easy going

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$Px$  : \_\_\_ $x$  is a person

## 2-Place Predicates

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Everyone is easygoing :

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Everyone is easygoing :

$$\forall x(\mathcal{F}x \rightarrow \mathcal{G}x)$$

## 2-Place Predicates

domain : everything in the office

|               |  |
|---------------|--|
| $j$ : Jim     | $Lxy$ : ___ $x$ likes ___ $y$          |
| $m$ : Michael | $Ex$ : ___ $x$ is easy going           |
| $p$ : Pam     | $Txy$ : ___ $x$ is taller than ___ $y$ |
| $s$ : Stanley | $Px$ : ___ $x$ is a person             |

Everyone is easygoing :

$$\forall x(Px \rightarrow Ex)$$

## 2-Place Predicates

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No one likes Michael :

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$Ex$  : \_\_\_ $x$  is easy going

$p$  : Pam

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

$Px$  : \_\_\_ $x$  is a person

No one likes Michael :

$$\forall x(\mathcal{F}x \rightarrow \neg \mathcal{G}x)$$

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$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

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$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

$Px$  : \_\_\_ $x$  is a person

No one likes Michael :

$$\forall x(Px \rightarrow \neg Lxm)$$

## 2-Place Predicates

domain : everything in the office

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$Ex$  : \_\_\_ $x$  is easy going

$p$  : Pam

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

$Px$  : \_\_\_ $x$  is a person

Michael likes everyone :

## 2-Place Predicates

domain : everything in the office

$j$  : Jim

$Lxy$  : \_\_\_ $x$  likes \_\_\_ $y$

$m$  : Michael

$Ey$  : \_\_\_ $x$  is easy going

$p$  : Pam

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

$Px$  : \_\_\_ $x$  is a person

Everyone is liked by Michael :

## 2-Place Predicates

domain : everything in the office

|               |  |
|---------------|--|
| $j$ : Jim     | $Lxy$ : ___ $x$ likes ___ $y$          |
| $m$ : Michael | $Ey$ : ___ $x$ is easy going           |
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Everyone is liked by Michael :

$$\forall x(\mathcal{F}x \rightarrow \mathcal{G}x)$$

## 2-Place Predicates

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Everyone is liked by Michael :

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## 2-Place Predicates

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| $s$ : Stanley | $Px$ : ___ $x$ is a person             |

Everyone is liked by Michael :

$$\forall x(Px \rightarrow Lmx)$$

## 2-Place Predicates

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$j$  : Jim

$Lxy$  : \_\_\_ $x$  likes \_\_\_ $y$

$m$  : Michael

$Ex$  : \_\_\_ $x$  is easy going

$p$  : Pam

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

$Px$  : \_\_\_ $x$  is a person

Stanley doesn't like anyone :

## 2-Place Predicates

domain : everything in the office

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$Lxy$  : \_\_\_ $x$  likes \_\_\_ $y$

$m$  : Michael

$Ex$  : \_\_\_ $x$  is easy going

$p$  : Pam

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

$Px$  : \_\_\_ $x$  is a person

No one is liked by Stanley :

## 2-Place Predicates

domain : everything in the office

$j$  : Jim

$Lxy$  : \_\_\_ $x$  likes \_\_\_ $y$

$m$  : Michael

$Ex$  : \_\_\_ $x$  is easy going

$p$  : Pam

$Txy$  : \_\_\_ $x$  is taller than \_\_\_ $y$

$s$  : Stanley

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No one is liked by Stanley :

$$\forall x(\mathcal{F}x \rightarrow \neg \mathcal{G}x)$$

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Someone likes Pam :

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| $j$ : Jim     | $Lxy$ : ___ $x$ likes ___ $y$          |
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| $s$ : Stanley | $Px$ : ___ $x$ is a person             |

Someone likes Pam :

$$\exists x(\mathcal{F}x \wedge \mathcal{G}x)$$



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Someone likes Pam :

$$\exists x(Px \wedge Lxp)$$

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Michael doesn't like anyone taller than him :

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No one taller than Michael is liked by Michael :

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No one taller than Michael is liked by Michael :

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No one taller than Michael is liked by Michael :

$$\forall x[(Px \wedge Txm) \rightarrow \neg \mathcal{L}x]$$

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No one taller than Michael is liked by Michael :

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Someone likes everyone :

$$\exists x [Px \wedge \forall y [Lyx]]$$

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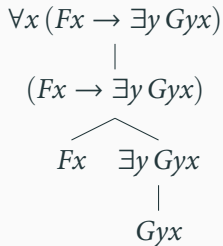
Someone likes everyone :

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# Syntax for PL

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PHIL 500





# Syntax for PL

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SYNTAX — { 1. Vocabulary  
2. Grammar  
SEMANTICS — 3. Meaning

# Syntax for PL

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## Vocabulary

# Vocabulary

The vocabulary of PL includes the following symbols:

1. for each  $N \geq 0$ ,  $N$ -place predicates (any capital letter—perhaps with subscripts)

$A, B, C, D, E, \dots, X, Y, Z$

$A_1, B_1, C_1, D_1, E_1, \dots, X_1, Y_1, Z_1$

$A_2, B_2, C_2, D_2, E_2, \dots, X_2, Y_2, Z_2$

$\vdots$

# Vocabulary

2. *names* (any lowercase letter between  $a$  and  $v$ —perhaps with subscripts)

$a, b, c, d, e, \dots, t, u, v$

$a_1, b_1, c_1, d_1, e_1, \dots, t_1, u_1, v_1$

$a_2, b_2, c_2, d_2, e_2, \dots, t_2, u_2, v_2$

$\vdots$

# Vocabulary

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$a_2, b_2, c_2, d_2, e_2, \dots, t_2, u_2, v_2$

$\vdots$

3. *variables* (lowercase  $w, x, y,$  and  $z$ —perhaps with subscripts)

$w, x, y, z$

$w_1, x_1, y_1, z_1$

$w_2, x_2, y_2, z_2$

$\vdots$

# Vocabulary

## 4. Logical operators

$\neg, \vee, \wedge, \rightarrow, \leftrightarrow, \exists, \forall$



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$(, )$

# Vocabulary

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## 5. parentheses

$(, )$

Nothing else is included in the vocabulary of PL.

- Let's call both names and variables *terms*. That is, both ' $a$ ' and ' $x$ ' are *terms* of PL.

# Syntax for PL

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## Grammar

- Any sequence of the symbols in the vocabulary of PL is an *expression* of PL.

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- All of the following are expressions of PL:

$$\forall x \neg ((\rightarrow \rightarrow \rightarrow a n v$$
$$P Q R S T \neg \neg$$
$$(\forall x F x a b \rightarrow \neg \exists y P y n s t)$$
$$N x y \vee \vee \neg \neg \exists x B x$$

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- All of the following are expressions of PL:

$Vx\neg((\rightarrow\rightarrow)anv)$

$PQRST\neg\neg$

$(\forall x Fx ab \rightarrow \neg\exists y Pynst)$

$Nxy \vee \vee\neg\neg\exists x Bx$



# Grammar: Atomic Sentences

- If  $\mathcal{R}$  is an  $N$ -place predicate and  $t_1, t_2, \dots, t_N$  are  $N$  terms, then

$$\mathcal{R}t_1t_2\dots t_N$$

is an atomic sentence.

## Grammar: Atomic sentences

- Let  $A$  be a 1-place predicate,  $B$  a 2-place predicate,  $C$  a 3-place predicate, and  $D$  a 4-place predicate

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- Let  $A$  be a 1-place predicate,  $B$  a 2-place predicate,  $C$  a 3-place predicate, and  $D$  a 4-place predicate
- Then, all of the following are atomic sentences of PL:

$Az$

$Aa$

$Bwg$

$Cxzt$

$Dcccc$

$Dxaxa$

## Grammar: Sentences

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- $\vee$ ) If ' $\mathcal{A}$ ' and ' $\mathcal{B}$ ' are sentences, then ' $(\mathcal{A} \vee \mathcal{B})$ ' is a sentence.
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- $\leftrightarrow$ ) If ' $\mathcal{A}$ ' and ' $\mathcal{B}$ ' are sentences, then ' $(\mathcal{A} \leftrightarrow \mathcal{B})$ ' is a sentence.
- $\forall$ ) If ' $\mathcal{A}$ ' is a sentence and ' $x$ ' is a variable, then ' $\forall x\mathcal{A}$ ' is a sentence.

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## Grammar: Sentences

- $\mathcal{R}$ ) Every atomic sentence is a sentence
- $\neg$ ) If ' $\mathcal{A}$ ' is a sentence, then ' $\neg\mathcal{A}$ ' is a sentence.
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- $-$ ) Nothing else is a sentence.

Note: none of ' $\mathcal{A}$ ', ' $\mathcal{B}$ ', ' $x$ ', or ' $t$ ' appear in the vocabulary of PL. They are not *themselves* sentences of PL. Rather, we are using them here as META-VARIABLES ranging over the expressions of PL.

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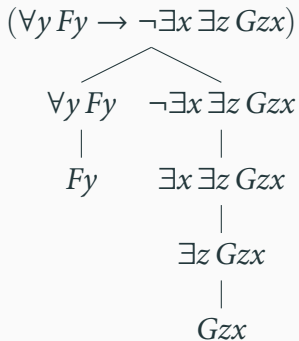
- Conventions:
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# Syntax Trees



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# Syntactic Structure

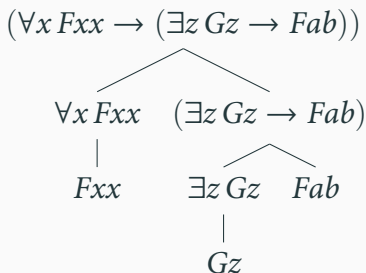
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# Subsentences

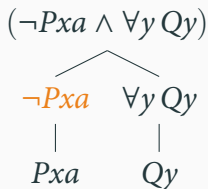
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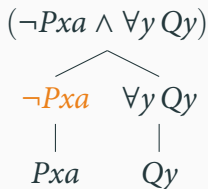
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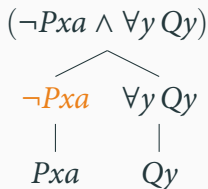
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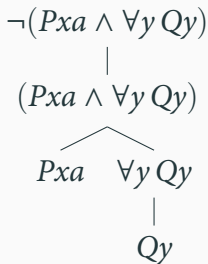
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# Main Operators

- The *main operator* in a (non-atomic) sentence is the operator which would be introduced *last*, if we were building the sentence up according to the rules for sentences.

# Main Operators

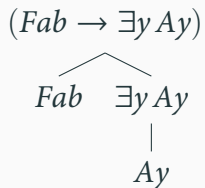
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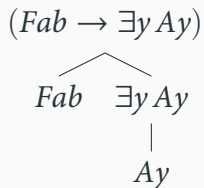
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main operator:  $\rightarrow$



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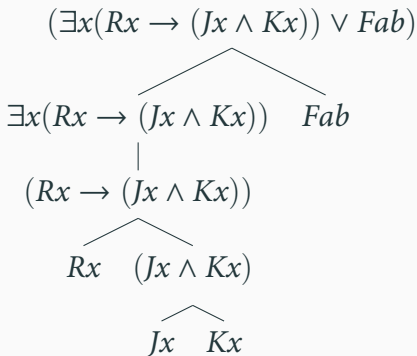
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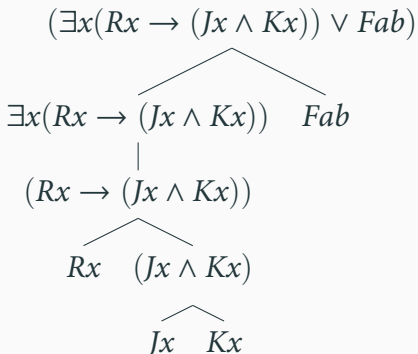
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# Main Operators

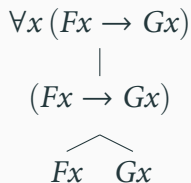
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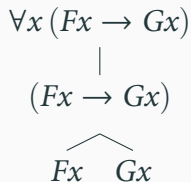
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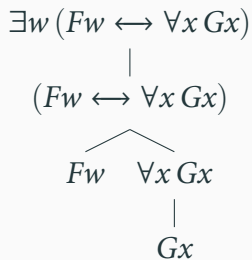
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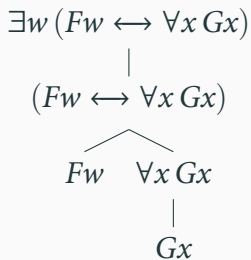
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- The *scope* of an operator (in a sentence) is the sub-sentence for which that operator is the main operator

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# Syntax for PL

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## Free and Bound Variables



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- The variables appearing in ' $\forall x \forall y Fxy$ ' are BOUND.
- In ' $\forall x Px \rightarrow Qx$ ', the first  $x$  is bound, whereas the second one is free.

## Free and Bound Variables

A variable  $x$  in a sentence of PL is **BOUND** if and only if it occurs within the scope of a quantifier,  $\forall x$  or  $\exists x$ , whose associated variable is  $x$ .

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A variable  $x$  in a sentence of PL is **FREE** if and only if it does *not* occur within the scope of a quantifier,  $\forall x$  or  $\exists x$ , whose associated variable is  $x$ .

- E.g.,

$$\forall x (\forall y Fy \rightarrow \exists z Gzx)$$

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In a sentence of the form  $\forall x \mathcal{A}$  or  $\exists x \mathcal{A}$ , the quantifier binds every *free* occurrence of  $x$  in  $\mathcal{A}$ . If an occurrence of  $x$  in  $\mathcal{A}$  is already bound, then the quantifier does not bind it.

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- E.g., in

$$\exists x \forall x Fx$$

the variable ' $x$ ' is bound by the *universal* quantifier ' $\forall x$ '. It is *not* bound by the existential quantifier ' $\exists x$ '.

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- When translating into PL, we want our translations to be *closed*.

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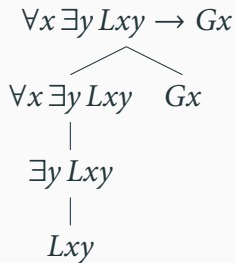
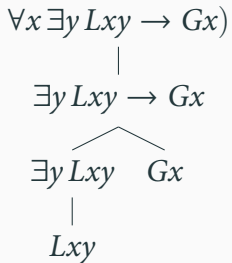
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# Syntax for PL

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## Important Syntactic Features in PL

# Parentheses



# Term Order

$$\forall x \exists y Lxy$$
$$|$$
$$\exists y Lxy$$
$$|$$
$$Lxy$$
$$\forall x \exists y Lyx$$
$$|$$
$$\exists y Lyx$$
$$|$$
$$Lyx$$

# Quantifier Order

$$\exists y \forall x Lxy$$
$$|$$
$$\forall x Lxy$$
$$|$$
$$Lxy$$
$$\forall x \exists y Lxy$$
$$|$$
$$\exists y Lxy$$
$$|$$
$$Lxy$$